Knowledge reduction of dynamic covering decision information systems with immigration of more objects

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Abstract. In practical situations, it is of interest to investigate computing approximations of sets as an important step of knowledge reduction of dynamic covering decision information systems. In this paper, we present incremental approaches to computing the type-1 and type-2 characteristic matrices of dynamic coverings whose cardinalities increase with immigration of more objects. We also present the incremental algorithms of computing the second and sixth lower and upper approximations of sets in dynamic covering approximation spaces.

Keywords: Boolean matrice; Characteristic matrice; Dynamic covering approximation space; Dynamic covering information system; Rough set

1 Introduction

Covering-based rough set theory [53], as a powerful mathematical tool for studying covering approximation spaces, has attracted a lot of attention of researchers in various fields of sciences. Especially, various kinds of approximation operators have been proposed for covering approximation spaces. Recently, Wang et al. [44] transformed the computation of approximations of a set into products of the characteristic matrices and the characteristic function of the set. However, it paid little attention to approaches to calculating the characteristic matrices. In practice, the covering approximation space varies with time due to the characteristics of data collection, and the non-incremental approach to constructing the characteristic matrices is often very costly or even intractable in dynamic covering approximation spaces. It is necessary to present effective approaches to computing characteristic matrices of dynamic coverings.

To the best of our knowledge, researchers [2, 3, 15, 26, 27, 54, 55] have focused on computing approximations of sets. For instance, Chen et al. [2, 3] constructed approximations of sets when coarsening or refining attribute values. Li et al. [15] computed approximations in dominance-based rough sets approach under the variation of attribute set. Luo et al. [26, 27] studied dynamic maintenance of approximations in set-valued ordered decision systems under the attribute generalization and the variation of object set.

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Zhang et al. [54,55] updated rough set approximations based on relation matrices and investigated neighborhood rough sets for dynamic data mining. These works demonstrate that incremental approaches are effective and efficient for computing approximations of sets. It motivates us to apply an incremental updating scheme to conduct approximations of sets by using characteristic matrices in dynamic covering approximation spaces, which will provide an effective approach to computing approximations of sets from the view of matrices.

The purpose of this paper is to compute approximations of sets by using incremental approaches in dynamic covering approximation spaces. First, we present incremental approaches to computing the type-1 and type-2 characteristic matrices in dynamic covering approximation spaces. We mainly focus on the situation: the variation of elements in coverings when adding and deleting objects. Furthermore, we provide incremental algorithms for constructing the second and sixth lower and upper approximations of sets based on the type-1 and type-2 characteristic matrices, respectively. We compare computation complexities of the incremental algorithms with those of non-incremental algorithms. Several examples are employed to illustrate that calculating approximations of sets is simplified greatly by utilizing the proposed approach.

The rest of this paper is organized as follows: Section 2 briefly reviews the basic concepts of covering-based rough set theory. In Section 3, we introduce incremental approaches to computing the type-1 and type-2 characteristic matrices of dynamic coverings whose cardinalities increase with immigration of more objects. In Section 4, we present incremental algorithms of calculating the second and sixth lower and upper approximations of sets by using the type-1 and type-2 characteristic matrices, respectively. We conclude the paper in Section 5.

2 Preliminaries

In this section, we briefly review some concepts of covering-based rough sets.

Definition 2.1 [53] Let U be a finite universe of discourse, and \mathscr{C} a family of subsets of U. Then \mathscr{C} is called a covering of U if none of elements of \mathscr{C} is empty and $\bigcup \{C | C \in \mathscr{C}\} = U$.

If \mathscr{C} is a covering of U, then (U,\mathscr{C}) is referred to as a covering approximation space.

Definition 2.2 [44] Let $U = \{x_1, x_2, ..., x_n\}$ be a finite universe, and $\mathcal{C} = \{C_1, C_2, ..., C_m\}$ a covering of U. For any $X \subseteq U$, the second, fifth and sixth upper and lower approximations of X with respect to \mathcal{C} are defined as

- $(1) SH_{\mathscr{C}}(X) = \bigcup \{C \in \mathscr{C} | C \cap X \neq \emptyset\}, SL_{\mathscr{C}}(X) = [SH_{\mathscr{C}}(X^c)]^c;$
- (2) $IH_{\mathscr{C}}(X) = \bigcup \{N(x)|N(x) \cap X \neq \emptyset, x \in U\}, IL_{\mathscr{C}}(X) = \bigcup \{N(x)|N(x) \subseteq X, x \in U\};$
- (3) $XH_{\mathscr{C}}(X) = \{x \in U | N(x) \cap X \neq \emptyset\}, XL_{\mathscr{C}}(X) = \{x \in U | N(x) \subseteq X\}, where N(x) = \bigcap \{C_i | x \in C_i \in \mathscr{C}\}.$

For simplicity, we omit \mathscr{C} in the following description of approximation operators.

Definition 2.3 [44] Let $U = \{x_1, x_2, ..., x_n\}$ be a finite universe, $\mathscr{C} = \{C_1, C_2, ..., C_m\}$ a family of subsets of U, and $M_{\mathscr{C}} = \{a_{ij}\}_{n \times m}$, where $a_{ij} = \{ \begin{array}{l} 1, & x_i \in C_j; \\ 0, & x_i \notin C_j. \end{array}$ Then $M_{\mathscr{C}}$ is called a matrice representation of

Accordingly, we have the characteristic function $X_X = \begin{bmatrix} a_1 & a_2 & \dots & a_n \end{bmatrix}^T$ for $X \subseteq U$, where $a_i = \begin{cases} 1, & x_i \in X; \\ 0, & x_i \notin X. \end{cases}$.

Definition 2.4 [44] Let \mathscr{C} be a covering of the universe U, $A = (a_{ij})_{n \times m}$ and $B = (b_{ij})_{m \times p}$ Boolean matrices, $A \odot B = (c_{ij})_{n \times p}$, where $c_{ij} = \bigwedge_{k=1}^{m} (b_{kj} - a_{ik} + 1)$. Then

- (1) $\Gamma(\mathscr{C}) = M_{\mathscr{C}} \cdot M_{\mathscr{C}}^T = (d_{ij})_{n \times n}$ is called the type-1 characteristic matrice of \mathscr{C} , where $d_{ij} = \bigvee_{k=1}^m (a_{ik} \cdot a_{jk})$, and $M_{\mathscr{C}} \cdot M_{\mathscr{C}}^T$ is the boolean product of $M_{\mathscr{C}}$ and its transpose $M_{\mathscr{C}}^T$;
 - (2) $\prod(\mathscr{C}) = M_{\mathscr{C}} \odot M_{\mathscr{C}}^T$ is referred to as the type-2 characteristic matrice of \mathscr{C} .

Wang et al. axiomatized two important types of covering approximation operators equivalently by using the type-1 and type-2 characteristic matrice of \mathscr{C} .

Definition 2.5 [44] Let $U = \{x_1, x_2, ..., x_n\}$ be a finite universe, $\mathcal{C} = \{C_1, C_2, ..., C_m\}$ a covering of U, and X_X the characteristic function of X in U. Then

$$(1)\ \mathcal{X}_{SH(X)} = \Gamma(\mathscr{C}) \cdot \mathcal{X}_X,\ \mathcal{X}_{SL(X)} = \Gamma(\mathscr{C}) \odot \mathcal{X}_X;\ (2)\ \mathcal{X}_{XH(X)} = \prod(\mathscr{C}) \cdot \mathcal{X}_X,\ \mathcal{X}_{XL(X)} = \prod(\mathscr{C}) \odot \mathcal{X}_X.$$

3 Update approximations of sets with immigration of more objects

In this section, we introduce incremental approaches to computing the second and sixth lower and upper approximation of sets with immigration of more objects.

Definition 3.1 Let (U, \mathcal{C}) and (U^+, \mathcal{C}^+) be covering approximation spaces, where $U = \{x_1, x_2, ..., x_n\}$, $U^+ = U \cup \{x_{n+1}, x_{n+2}, ..., x_{n+t}\} (t \ge 2)$, $\mathcal{C} = \{C_1, C_2, ..., C_m\}$, $\mathcal{C}^+ = \{C_1^+, C_2^+, ..., C_m^+, C_{m+1}^+, C_{m+2}^+, ..., C_{m+l}^+\} (l \ge 2)$, where $C_i^+ = C_i \cup \Delta C_i$ or C_i $(1 \le i \le m)$, $\Delta C_i \subseteq \{x_{n+1}, x_{n+2}, ..., x_{n+t}\}$, and $\{x_{n+1}, x_{n+2}, ..., x_{n+t}\} \subseteq \{C_{m+j}^+ | 1 \le j \le l\}$. Then (U^+, \mathcal{C}^+) is called a dynamic covering approximation space.

By Definition 3.1, we refer \mathscr{C}^+ to as a dynamic covering. Although there are several types of coverings when adding objects, we only discuss this type of dynamic coverings for simplicity in this work.

In what follows, we discuss how to construct $\Gamma(\mathscr{C}^+)$ based on $\Gamma(\mathscr{C})$. For convenience, we denote $M_{\mathscr{C}} = (a_{ij})_{n \times m}$, $M_{\mathscr{C}^+} = (a_{ij})_{(n+t)\times(m+l)}$, $\Gamma(\mathscr{C}) = (b_{ij})_{n \times n}$ and $\Gamma(\mathscr{C}^+) = (c_{ij})_{(n+t)\times(n+t)}$.

Theorem 3.2 Let (U^+, \mathcal{C}^+) be a dynamic covering approximation space of (U, \mathcal{C}) , $\Gamma(\mathcal{C})$ and $\Gamma(\mathcal{C}^+)$ the type-1 characteristic matrices of \mathcal{C} and \mathcal{C}^+ , respectively. Then

$$\Gamma(\mathscr{C}^+) = \left[\begin{array}{cc} \Gamma(\mathscr{C}) & 0 \\ 0 & 0 \end{array} \right] \bigvee \left[\begin{array}{cc} \triangle_1(\Gamma(\mathscr{C})) & (\triangle_2(\Gamma(\mathscr{C})))^T \\ \triangle_2(\Gamma(\mathscr{C})) & \triangle_3(\Gamma(\mathscr{C})) \end{array} \right],$$

where

$$\Delta_{1}(\Gamma(\mathscr{C})) \ = \ \begin{bmatrix} a_{1(m+1)} & a_{2(m+1)} & \dots & a_{n(m+1)} \\ a_{1(m+2)} & a_{2(m+2)} & \dots & a_{n(m+2)} \\ \vdots & \vdots & \ddots & \vdots \\ a_{1(m+l)} & a_{2(m+l)} & \dots & a_{n(m+2)} \end{bmatrix}^{T} \begin{bmatrix} a_{1(m+1)} & a_{2(m+1)} & \dots & a_{n(m+1)} \\ a_{1(m+2)} & a_{2(m+2)} & \dots & a_{n(m+2)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{1(m+l)} & a_{2(m+l)} & \dots & a_{n(m+2)} \end{bmatrix}; \\ \Delta_{2}(\Gamma(\mathscr{C})) \ = \ \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{n(m+l)} \\ a_{(n+2)1} & a_{(n+2)2} & \dots & a_{(n+2)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+l)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix} \cdot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1(m+l)} \\ a_{21} & a_{22} & \dots & a_{2(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{n(m+l)} \end{bmatrix}^{T} ; \\ \Delta_{3}(\Gamma(\mathscr{C})) \ = \ \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+l)(m+l)} \\ a_{(n+2)1} & a_{(n+2)2} & \dots & a_{(n+2)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+l)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix} \cdot \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)(m+l)} \\ a_{(n+2)1} & a_{(n+2)2} & \dots & a_{(n+2)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+l)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix}^{T}$$

Proof. By Definition 3.1, we get $\Gamma(\mathscr{C})$ and $\Gamma(\mathscr{C}^+)$ as follows:

$$\Gamma(\mathcal{C}) = M_{\mathcal{C}} \cdot M_{\mathcal{C}}^{T}$$

$$= \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} \\ a_{21} & a_{22} & \dots & a_{2m} \\ \vdots & \vdots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nm} \end{bmatrix} \cdot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} \\ a_{21} & a_{22} & \dots & a_{2m} \\ \vdots & \vdots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nm} \end{bmatrix}^{T}$$

$$= \begin{bmatrix} b_{11} & b_{12} & \dots & b_{1n} \\ b_{21} & b_{22} & \dots & b_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ b_{n1} & b_{n2} & \dots & b_{nn} \end{bmatrix};$$

$$\Gamma(\mathscr{C}^+) \ = \ M_{\mathscr{C}^+} \cdot M_{\mathscr{C}^+}^+ \\ = \ \begin{bmatrix} a_{11} & a_{12} & & a_{1m} & a_{1(m+1)} & & a_{1(m+l)} \\ a_{21} & a_{22} & & a_{2m} & a_{2(m+1)} & & a_{2(m+l)} \\ \vdots & \vdots & & \vdots & & \vdots & & \vdots \\ \vdots & \vdots & & \ddots & & \vdots & & \vdots \\ a_{n1} & a_{n2} & & a_{nm} & a_{n(m+1)} & & a_{n(m+l)} \\ a_{(n+1)1} & a_{(n+1)2} & & a_{(n+1)m} & a_{(n+1)(m+1)} & & & a_{(n+1)(m+l)} \\ \vdots & & & & & \vdots & & \vdots \\ a_{(n+t)1} & a_{(n+t)2} & & a_{(n+t)m} & a_{(n+t)(m+1)} & & & a_{(n+t)(m+l)} \end{bmatrix}^{1} \\ \begin{bmatrix} a_{11} & a_{12} & & a_{1m} & a_{1(m+1)} & & a_{1(m+l)} \\ a_{21} & a_{22} & & a_{2m} & a_{2(m+1)} & & a_{2(m+l)} \\ \vdots & & & & & \vdots \\ \vdots & & & & & \vdots \\ a_{n1} & a_{n2} & & a_{nm} & a_{n(m+1)} & & a_{n(m+l)} \\ a_{(n+1)1} & a_{(n+1)2} & & a_{(n+1)m} & a_{(n+1)(m+1)} & & a_{(n+1)(m+l)} \\ \vdots & & & & & & \vdots \\ \vdots & & & & & & \vdots \\ a_{(n+t)1} & a_{(n+t)2} & & a_{(n+t)m} & a_{(n+t)(m+1)} & & a_{(n+t)(m+l)} \\ \end{bmatrix}^{1} \\ \begin{bmatrix} c_{11} & c_{12} & & c_{1n} & c_{1(n+1)} & & c_{1(n+t)} \\ c_{21} & c_{22} & & c_{2n} & c_{2(n+1)} & & c_{2(n+t)} \\ \vdots & & & & & \vdots \\ \vdots & & & & & & \vdots \\ c_{(n+1)1} & c_{(n+1)2} & & c_{(n+t)n} & c_{(n+t)(n+1)} & & c_{(n+t)(n+t)} \end{bmatrix}^{1} \\ \end{bmatrix}^{1} \\ \begin{bmatrix} c_{11} & c_{12} & & c_{mn} & c_{n(n+1)} & & c_{n(n+t)} \\ c_{21} & c_{22} & & c_{2n} & c_{2(n+1)} & & c_{2(n+t)} \\ \vdots & & & & & \vdots \\ c_{(n+t)1} & c_{(n+t)2} & & c_{(n+t)n} & c_{(n+t)(n+1)} & & c_{(n+t)(n+t)} \end{bmatrix}^{1} \\ \end{bmatrix}^{1} \\ \begin{bmatrix} c_{11} & c_{12} & & c_{mn} & c_{n(n+1)} & & c_{n(n+t)} \\ c_{21} & c_{22} & & c_{2n} & c_{2(n+1)} & & c_{2(n+t)} \\ \vdots & \vdots & & \vdots & & \vdots \\ c_{(n+t)1} & c_{(n+t)2} & & c_{(n+t)n} & c_{(n+t)(n+1)} & & c_{(n+t)(n+t)} \end{bmatrix}^{1}$$

In the sense of the type-1 characteristic matrice of \mathcal{C}^+ , we have

$$c_{11} = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} & a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T}$$

$$\cdot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} & a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T}$$

$$= \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} \end{bmatrix} \cdot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} \end{bmatrix}^{T}$$

$$\vee \begin{bmatrix} a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix} \cdot \begin{bmatrix} a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T}$$

$$= b_{11} \vee \begin{bmatrix} a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix} \cdot \begin{bmatrix} a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T};$$

$$c_{(n+1)1} = \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}$$

$$\cdot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} & a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T}$$

$$= 0 \vee \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} & a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T};$$

$$c_{(n+1)(n+1)} = \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}^{T}$$

$$\cdot \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}^{T}$$

$$= 0 \lor \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}^{T}$$

$$\cdot \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}^{T}.$$

Since $c_{11} \in \Delta_1(\Gamma(\mathscr{C}))$, $c_{(n+1)1} \in \Delta_2(\Gamma(\mathscr{C}))$ and $c_{(n+1)(n+1)} \in \Delta_3(\Gamma(\mathscr{C}))$, we can compute other elements of $\Delta_1(\Gamma(\mathscr{C}))$, $\Delta_2(\Gamma(\mathscr{C}))$ and $\Delta_3(\Gamma(\mathscr{C}))$ similarly. Thus, to obtain $\Gamma(\mathscr{C}^+)$, we only need to compute $\Delta_1(\Gamma(\mathscr{C}))$, $\Delta_2(\Gamma(\mathscr{C}))$ and $\Delta_3(\Gamma(\mathscr{C}))$ on the basis of $\Gamma(\mathscr{C})$ as follows:

$$\Delta_{1}(\Gamma(\mathscr{C})) = \begin{bmatrix} a_{1(m+1)} & a_{2(m+1)} & \dots & a_{n(m+1)} \\ a_{1(m+2)} & a_{2(m+2)} & \dots & a_{n(m+2)} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ a_{1(m+l)} & a_{2(m+l)} & \dots & a_{n(m+l)} \end{bmatrix}^{T} \begin{bmatrix} a_{1(m+1)} & a_{2(m+1)} & \dots & a_{n(m+1)} \\ a_{1(m+2)} & a_{2(m+2)} & \dots & a_{n(m+2)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{1(m+l)} & a_{2(m+l)} & \dots & a_{n(m+l)} \end{bmatrix}^{T} \\ \Delta_{2}(\Gamma(\mathscr{C})) = \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)(m+l)} \\ a_{(n+2)1} & a_{(n+2)2} & \dots & a_{(n+1)(m+l)} \\ \vdots & \vdots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+l)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix} \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1(m+l)} \\ a_{21} & a_{22} & \dots & a_{2(m+l)} \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{n(m+l)} \end{bmatrix}^{T} \\ \vdots & \vdots & \ddots & \vdots \\ a_{(n+2)1} & a_{(n+2)2} & \dots & a_{(n+2)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+l)2} & \dots & a_{(n+l)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+l)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix}^{T} \\ \vdots & \vdots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots &$$

Therefore, we have

$$\Gamma(\mathscr{C}^+) = \begin{bmatrix} \Gamma(\mathscr{C}) & 0 \\ 0 & 0 \end{bmatrix} \bigvee \begin{bmatrix} \Delta_1(\Gamma(\mathscr{C})) & (\Delta_2(\Gamma(\mathscr{C})))^T \\ \Delta_2(\Gamma(\mathscr{C})) & \Delta_3(\Gamma(\mathscr{C})) \end{bmatrix}.$$

Example 3.3 Let $U = \{x_1, x_2, x_3, x_4\}$, $U^+ = U \cup \{x_5, x_6\}$, $\mathscr{C} = \{C_1, C_2, C_3\}$, $\mathscr{C}^+ = \{C_1^+, C_2^+, C_3^+, C_4^+, C_5^+\}$, where $C_1 = \{x_1, x_4\}$, $C_2 = \{x_1, x_2, x_4\}$, $C_3 = \{x_3, x_4\}$, $C_1^+ = \{x_1, x_4, x_5\}$, $C_2^+ = \{x_1, x_2, x_4, x_5\}$, $C_3^+ = \{x_3, x_4\}$, $C_4^+ = \{x_3, x_5, x_6\}$, $C_5^+ = \{x_1, x_6\}$, and $X = \{x_3, x_4, x_5\}$. By Definition 3.1, we first have that

$$\Gamma(\mathscr{C}) \ = \ M_{\mathscr{C}} \cdot M_{\mathscr{C}}^T = \left[\begin{array}{ccc} 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 1 \end{array} \right] \cdot \left[\begin{array}{ccc} 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 1 \end{array} \right]^T = \left[\begin{array}{ccc} 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{array} \right].$$

Second, by Theorem 3.2, we get that

$$\triangle_{1}(\Gamma(\mathscr{C})) = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}^{T} \cdot \begin{bmatrix} 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix};$$

$$\Delta_{2}(\Gamma(\mathscr{C})) = \begin{bmatrix} 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \end{bmatrix} \cdot \begin{bmatrix} 1 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 \end{bmatrix}^{T} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \end{bmatrix};$$

$$\Delta_{3}(\Gamma(\mathscr{C})) = \begin{bmatrix} 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \end{bmatrix} \cdot \begin{bmatrix} 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \end{bmatrix}^{T} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}.$$

Thus, we obtain that

By Definition 2.5, we have that

Therefore, $SH(X) = \{x_1, x_2, x_3, x_4, x_5, x_6\}$ and $SL(X) = \{x_3\}$.

In Example 3.3, we only need to calculate elements in $\triangle_1(\Gamma(\mathscr{C}))$, $\triangle_2(\Gamma(\mathscr{C}))$ and $\triangle_3(\Gamma(\mathscr{C}))$ by Theorem 3.2. Thereby, the incremental algorithm is effective to compute the second lower and upper approximations of sets.

In practical situations, there exists a need to construct the type-2 characteristic matrices of dynamic coverings for computing the sixth lower and upper approximations of sets. Subsequently, we construct $\prod(\mathscr{C}^+)$ based on $\prod(\mathscr{C})$. For convenience, we denote $\prod(\mathscr{C}) = (d_{ij})_{n \times n}$ and $\prod(\mathscr{C}^+) = (e_{ij})_{(n+t) \times (n+t)}$.

Theorem 3.4 Let (U^+, \mathscr{C}^+) be a dynamic covering approximation space of (U, \mathscr{C}) , $\Pi(\mathscr{C})$ and $\Pi(\mathscr{C}^+)$ the

type-2 characteristic matrices of $\mathscr C$ and $\mathscr C^+$, respectively. Then

$$\prod(\mathscr{C}^+) = \left[\begin{array}{cc} \prod(\mathscr{C}) & 1 \\ 1 & 1 \end{array} \right] \bigwedge \left[\begin{array}{cc} \triangle_1(\prod(\mathscr{C})) & \triangle_3(\prod(\mathscr{C})) \\ \triangle_2(\prod(\mathscr{C})) & \triangle_4(\prod(\mathscr{C})) \end{array} \right],$$

where

$$\Delta_{1}(\prod(\mathscr{C})) \ = \ \begin{bmatrix} a_{1(m+1)} & a_{2(m+1)} & \dots & a_{n(m+1)} \\ a_{1(m+2)} & a_{2(m+2)} & \dots & a_{n(m+2)} \\ \vdots & \vdots & \ddots & \vdots \\ a_{1(m+l)} & a_{2(m+l)} & \dots & a_{n(m+l)} \end{bmatrix}^{T} \odot \begin{bmatrix} a_{1(m+1)} & a_{2(m+1)} & \dots & a_{n(m+1)} \\ a_{1(m+2)} & a_{2(m+2)} & \dots & a_{n(m+2)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{1(m+l)} & a_{2(m+l)} & \dots & a_{n(m+l)} \end{bmatrix};$$

$$\Delta_{2}(\prod(\mathscr{C})) \ = \ \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)(m+l)} \\ a_{(n+2)1} & a_{(n+2)2} & \dots & a_{(n+2)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+l)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix} \odot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1(m+l)} \\ a_{21} & a_{22} & \dots & a_{2(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{n(m+l)} \end{bmatrix}^{T};$$

$$\Delta_{3}(\prod(\mathscr{C})) \ = \ \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1(m+l)} \\ a_{21} & a_{22} & \dots & a_{2(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{n(m+l)} \end{bmatrix} \odot \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)(m+l)} \\ a_{(n+2)1} & a_{(n+2)2} & \dots & a_{(n+1)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+2)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix}^{T};$$

$$\Delta_{4}(\prod(\mathscr{C})) \ = \ \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)(m+l)} \\ a_{(n+2)1} & a_{(n+2)2} & \dots & a_{(n+1)(m+l)} \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+l)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix}^{T}.$$

Proof. By Definition 3.1, we have $\prod(\mathscr{C})$ and $\prod(\mathscr{C}^+)$ as follows:

$$\prod(\mathscr{C}) = M_{\mathscr{C}} \odot M_{\mathscr{C}}^{T} \\
= \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} \\ a_{21} & a_{22} & \dots & a_{2m} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nm} \end{bmatrix} \odot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} \\ a_{21} & a_{22} & \dots & a_{2m} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nm} \end{bmatrix}^{T} \\
= \begin{bmatrix} d_{11} & d_{12} & \dots & d_{1n} \\ d_{21} & d_{22} & \dots & d_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \vdots \\ d_{n1} & d_{n2} & \dots & d_{nn} \end{bmatrix};$$

In the sense of the type-2 characteristic matrice of \mathscr{C}^+ , we have

$$e_{11} = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} & a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T}$$

$$\odot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} & a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T}$$

$$= \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} \end{bmatrix} \odot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} \end{bmatrix}^{T}$$

$$\wedge \begin{bmatrix} a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix} \odot \begin{bmatrix} a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T}$$

$$= d_{11} \wedge \begin{bmatrix} a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix} \odot \begin{bmatrix} a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T};$$

$$e_{(n+1)1} = \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}$$

$$\odot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} & a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T}$$

$$= 1 \wedge \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{1m} & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}$$

$$\odot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} & a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}^{T};$$

 $e_{(n+t)2}$. . . $e_{(n+t)n}$ $e_{(n+t)(n+1)}$. . .

$$e_{1(n+1)} = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} & a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}$$

$$\odot \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}^T$$

$$= 1 \land \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} & a_{1(m+1)} & \dots & a_{1(m+l)} \end{bmatrix}$$

$$\odot \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}^T;$$

$$e_{(n+1)(n+1)} = \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}^T;$$

$$= 1 \land \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}^T$$

$$= 1 \land \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}^T.$$

$$\odot \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)m} & a_{(n+1)(m+1)} & \dots & a_{(n+1)(m+l)} \end{bmatrix}^T.$$

Since $e_{11} \in \triangle_1(\prod(\mathscr{C}))$, $e_{(n+1)1} \in \triangle_2(\prod(\mathscr{C}))$, $e_{1(n+1)} \in \triangle_3(\prod(\mathscr{C}))$ and $e_{(n+1)(n+1)} \in \triangle_3(\prod(\mathscr{C}))$, we can compute other elements of $\triangle_1(\prod(\mathscr{C}))$, $\triangle_2(\prod(\mathscr{C}))$, $\triangle_3(\prod(\mathscr{C}))$ and $\triangle_4(\prod(\mathscr{C}))$ similarly. Thus, to compute $\prod(\mathscr{C}^+)$ on the basis of $\prod(\mathscr{C})$, we only need to compute $\triangle_1(\prod(\mathscr{C}))$, $\triangle_2(\prod(\mathscr{C}))$, $\triangle_3(\prod(\mathscr{C}))$ and $\triangle_4(\prod(\mathscr{C}))$ as follows:

$$\Delta_{1}(\prod(\mathscr{C})) \ = \ \begin{bmatrix} a_{1(m+1)} & a_{2(m+1)} & \dots & a_{n(m+1)} \\ a_{1(m+2)} & a_{2(m+2)} & \dots & a_{n(m+2)} \\ \vdots & \vdots & \ddots & \vdots \\ a_{1(m+l)} & a_{2(m+l)} & \dots & a_{n(m+2)} \end{bmatrix}^{T} \odot \begin{bmatrix} a_{1(m+1)} & a_{2(m+1)} & \dots & a_{n(m+1)} \\ a_{1(m+2)} & a_{2(m+2)} & \dots & a_{n(m+2)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{1(m+l)} & a_{2(m+2)} & \dots & a_{n(m+2)} \end{bmatrix};$$

$$\Delta_{2}(\prod(\mathscr{C})) \ = \ \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+1)(m+l)} \\ a_{(n+2)1} & a_{(n+2)2} & \dots & a_{(n+2)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+l)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix}^{T} \odot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1(m+l)} \\ a_{21} & a_{22} & \dots & a_{2(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{n(m+l)} \end{bmatrix}^{T} \odot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1(m+l)} \\ a_{21} & a_{22} & \dots & a_{2(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{n(m+l)} \end{bmatrix}^{T} \odot \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1(m+l)} \\ a_{21} & a_{22} & \dots & a_{2(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+2)2} & \dots & a_{(n+l)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+l)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix}^{T} \odot \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+l)(m+l)} \\ a_{(n+2)1} & a_{(n+2)2} & \dots & a_{(n+l)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+2)2} & \dots & a_{(n+l)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+2)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix}^{T} \odot \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+l)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+2)2} & \dots & a_{(n+l)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+2)2} & \dots & a_{(n+l)(m+l)} \end{bmatrix}^{T} \odot \begin{bmatrix} a_{(n+1)1} & a_{(n+1)2} & \dots & a_{(n+l)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{(n+l)1} & a_{(n+2)2} & \dots & a_{(n+l)(m+l)} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \vdots & \vdots$$

Therefore, we have

$$\prod(\mathscr{C}^+) = \left[\begin{array}{cc} \prod(\mathscr{C}) & 1 \\ 1 & 1 \end{array} \right] \bigwedge \left[\begin{array}{cc} \triangle_1(\prod(\mathscr{C})) & \triangle_3(\prod(\mathscr{C})) \\ \triangle_2(\prod(\mathscr{C})) & \triangle_4(\prod(\mathscr{C})) \end{array} \right].$$

The following example illustrates that how to compute the sixth lower and upper approximations of set by using the incremental algorithm.

Example 3.5 (Continuation of Example 3.3) We obtain that

$$\prod(\mathscr{C}) = M_{\mathscr{C}} \odot M_{\mathscr{C}}^T = \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 1 \end{bmatrix} \odot \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 1 \end{bmatrix}^T = \begin{bmatrix} 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{bmatrix}.$$

By Theorem 3.4, we have that

$$\Delta_{1}(\prod(\mathscr{C})) = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}^{T} \odot \begin{bmatrix} 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 \end{bmatrix};$$

$$\Delta_{2}(\prod(\mathscr{C})) = \begin{bmatrix} 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \end{bmatrix} \odot \begin{bmatrix} 1 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 \end{bmatrix}^{T} = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix};$$

$$\Delta_{3}(\prod(\mathscr{C})) = \begin{bmatrix} 1 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 \end{bmatrix}^{T} \odot \begin{bmatrix} 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix};$$

$$\Delta_{4}(\prod(\mathscr{C})) = \begin{bmatrix} 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \end{bmatrix} \odot \begin{bmatrix} 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \end{bmatrix}^{T} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}.$$

Thus, we have

$$\prod(\mathscr{C}^{+}) \ = \ \left[\begin{array}{ccc} \prod(\mathscr{C}) & 1 \\ 1 & 1 \end{array}\right] \bigwedge \left[\begin{array}{ccc} \triangle_{1}(\prod(\mathscr{C})) & \triangle_{3}(\prod(\mathscr{C})) \\ \triangle_{2}(\prod(\mathscr{C})) & \triangle_{4}(\prod(\mathscr{C})) \end{array}\right] = \left[\begin{array}{cccc} 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{array}\right].$$

By Definition 2.5, we obtain

$$\mathcal{X}_{XH(X)} = \prod (\mathscr{C}^{+}) \cdot \mathcal{X}_{X} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix} \cdot \begin{bmatrix} 0 \\ 0 \\ 1 \\ 1 \\ 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 1 & 1 & 1 & 0 \end{bmatrix}^{T};$$

$$\mathcal{X}_{XL(X)} = \prod (\mathscr{C}^{+}) \odot \mathcal{X}_{X} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix} \odot \begin{bmatrix} 0 \\ 0 \\ 1 \\ 1 \\ 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 0 & 0 & 1 & 1 & 1 & 0 \end{bmatrix}^{T}.$$

Therefore, $XH(X) = \{x_2, x_3, x_4, x_5\}$ and $XL(X) = \{x_3, x_4, x_5\}$.

In Example 3.5, we need to compute all elements in $\prod(\mathscr{C}^+)$ for constructing approximations of sets by Definition 3.1. By Theorem 3.4, we only need to calculate elements in $\triangle_1(\prod(\mathscr{C}))$, $\triangle_2(\prod(\mathscr{C}))$, $\triangle_3(\prod(\mathscr{C}))$ and $\triangle_4(\prod(\mathscr{C}))$. Thereby, the incremental algorithm is more effective to compute approximations of sets.

Non-incremental and incremental algorithms for computing the second and sixth lower and upper approximations of sets

In this section, we show non-incremental and incremental algorithms of computing the second lower and upper approximations of sets.

Algorithm 4.1 (Non-incremental algorithm of computing $SH_{\mathscr{C}^+}(X^+)$ and $SL_{\mathscr{C}^+}(X^+)(NCS)$)

Step 1: Input
$$(U^+, \mathcal{C}^+)$$
 and $X^+ \subseteq U^+$;

Step 2: Construct
$$M_{\mathcal{C}^+}$$
 and $\Gamma(\mathcal{C}^+) = M_{\mathcal{C}^+} \cdot M_{\mathcal{C}^+}^T$;

Step 3: Compute
$$X_{SH_{\mathscr{L}^+}(X^+)} = \Gamma(\mathscr{C}^+) \cdot X_{X^+}$$
 and $X_{SL_{\mathscr{L}^+}(X^+)} = \Gamma(\mathscr{C}^+) \odot X_{X^+}$;

Step 4: Output
$$SH_{\mathcal{C}^+}(X^+)$$
 and $SL_{\mathcal{C}^+}(X^+)$.

Algorithm 4.2 (Incremental algorithm of computing $SH_{\mathcal{C}^+}(X^+)$ and $SL_{\mathcal{C}^+}(X^+)$ (ICS))

Step 1: Input
$$(U, \mathcal{C})$$
, (U^+, \mathcal{C}^+) and $X \subseteq U^+$;

Step 2: Calculate
$$\Gamma(\mathcal{C}) = M_{\mathcal{C}} \cdot M_{\mathcal{C}}^T$$
, where $M_{\mathcal{C}} = (a_{ij})_{n \times m}$;

Step 3: Compute
$$\triangle_1(\Gamma(\mathscr{C}))$$
 and $\triangle_2(\Gamma(\mathscr{C}))$ and $\triangle_3(\Gamma(\mathscr{C}))$;

Step 4: Construct $\Gamma(\mathcal{C}^+)$, where

$$\Gamma(\mathscr{C}^{+}) = (c_{ij})_{(n+1)(n+1)} = \begin{bmatrix} \Gamma(\mathscr{C}) & 0 \\ 0 & 0 \end{bmatrix} \bigvee \begin{bmatrix} \Delta_{1}(\Gamma(\mathscr{C})) & (\Delta_{2}(\Gamma(\mathscr{C})))^{T} \\ \Delta_{2}(\Gamma(\mathscr{C})) & \Delta_{3}(\Gamma(\mathscr{C})) \end{bmatrix};$$

Step 5: Obtain $X_{SH(X)}$ and $X_{SL(X)}$, where

$$X_{SH(X)} = \Gamma(\mathscr{C}^+) \cdot X_X; X_{SL(X)} = \Gamma(\mathscr{C}^+) \odot X_X.$$

Subsequently, we present non-incremental and incremental algorithms of computing the sixth lower and upper approximations of sets.

Algorithm 4.3 (Non-incremental algorithm of computing $XH_{\mathscr{C}^+}(X^+)$ and $XL_{\mathscr{C}^+}(X^+)(\mathbf{NCX})$)

Step 1: Input
$$(U^+, \mathcal{C}^+)$$
 and $X^+ \subseteq U^+$;

Step 2: Construct
$$M_{\mathcal{C}^+}$$
 and $\prod(\mathcal{C}^+) = M_{\mathcal{C}^+} \cdot M_{\mathcal{C}^+}^T$;

Step 3: Compute
$$X_{XH_{\mathscr{L}^+}(X^+)} = \prod (\mathscr{C}^+) \cdot X_{X^+}$$
 and $X_{XL_{\mathscr{L}^+}(X^+)} = \prod (\mathscr{C}^+) \odot X_{X^+}$;

Step 4: Output
$$XH_{\mathcal{C}^+}(X^+)$$
 and $XL_{\mathcal{C}^+}(X^+)$.

Algorithm 4.4 (Incremental algorithm of computing $XH_{\mathcal{C}^+}(X^+)$ and $XL_{\mathcal{C}^+}(X^+)(\mathbf{ICX})$)

Step 1: Input
$$(U, \mathcal{C})$$
, (U^+, \mathcal{C}^+) and $X \subseteq U^+$;

Step 2: Construct
$$\prod(\mathscr{C})$$
, where $\prod(\mathscr{C}) = M_{\mathscr{C}} \odot M_{\mathscr{C}}^T$;

Step 3: Compute
$$\wedge_1(\Pi(\mathcal{E}))$$
 and $\wedge_2(\Pi(\mathcal{E}))$ $\wedge_2(\Pi(\mathcal{E}))$ and $\wedge_4(\Pi(\mathcal{E}))$:

Step 3: Compute
$$\triangle_1(\prod(\mathscr{C}))$$
 and $\triangle_2(\prod(\mathscr{C}))$, $\triangle_3(\prod(\mathscr{C}))$ and $\triangle_4(\prod(\mathscr{C}))$;
Step 4: Calculate $\prod(\mathscr{C}^+)$, where $\prod(\mathscr{C}^+) = \begin{bmatrix} \prod(\mathscr{C}) & 1 \\ 1 & 1 \end{bmatrix} \wedge \begin{bmatrix} \triangle_1(\prod(\mathscr{C})) & \triangle_3(\prod(\mathscr{C})) \\ \triangle_2(\prod(\mathscr{C})) & \triangle_4(\prod(\mathscr{C})) \end{bmatrix}$;

Step 5: Get $X_{XH(X)}$ and $X_{XL(X)}$, where

$$\mathcal{X}_{XH(X)} = \prod (\mathcal{C}^+) \cdot \mathcal{X}_X; \mathcal{X}_{XL(X)} = \prod (\mathcal{C}^+) \odot \mathcal{X}_X.$$

5 Conclusions

In this paper, we have provided effective approaches to constructing approximations of concepts in dynamic covering approximation spaces. Concretely, we have constructed type-1 and type-2 characteristic matrices of coverings with the incremental approaches. Incremental algorithms have been presented for computing the second and sixth lower and upper approximations of sets. Several examples have been employed to illustrate that computing approximations of sets could be reduced greatly by using the incremental approaches.

In the future, we will propose more effective approaches to constructing the type-1 and type-2 characteristic matrices of coverings. Additionally, we will focus on the development of effective approaches for knowledge discovery in dynamic covering approximation spaces.

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